Cyber-Physical Systems Modeling and Simulation with Modelica

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Abstract

This paper introduces the area of Cyber-Physical Systems (CPS) and describes the relation to Modelica and Modelica-based tools. Special aspects of CPS applications that should make Modelica well suited for their modeling and simulation are highlighted.

Recent Modelica developments facilitating integrated model-based system development applicable to CPS are presented. Especially, it is shown how detailed timing simulations, involving both real-time task scheduling and network communication, are realized in Modelica. A robot example is used to demonstrate the new CPS simulation features.

Keywords: Cyber-physical systems, Modelica, Model-based development, Timing simulation.

1 Introduction

During the last few years the field of Cyber-Physical Systems (CPS) has emerged, mainly in the US, as one of the most important academic and industrial research topics for the future. CPS focus on the interaction between computing/communication and the physical world. The following definition of CPS is taken from Wikipedia [1]

"A cyber-physical system (CPS) is a system featuring a tight combination of, and coordination between, the system's **computational** and **physical** elements."

Embedded computing devices controlling physical processes are today found in many diverse application areas, including automotive, aerospace, energy, and telecommunications. These traditional embedded systems can be viewed as a sub-set of the wider CPS definition. Typical CPS applications tend to be large-scale distributed systems built up from networks of computing/communication devices and possibly also distributed physical plants, such as in power grids. The complexity and scale inherent to CPS should make Modelica well suited for modeling and simulation of these systems.

The US National Science Foundation (NSF) has identified cyber-physical systems as a key area of research and several workshops have been organized on various aspects of CPS. The international multiconference, CPSWeek, has been organized since 2008 and *The Cyber-Physical Systems Summit* was held in conjunction with the first CPSWeek in St Louis, Missouri. This resulted in a report [2] that outlines the CPS field, including important applications and future research challenges. One of the observations in the report is that traditionally

"... the fields of computer science and control theory have remained largely separate, both technically and culturally."

To be able to build complex applications that successfully combine computational elements with elements from the physical world, it is important to have a new, integrated approach to the design. Modelica and Modelica-based tools should comprise good environments for such integrated model-based development. The report goes on to state that

"On the one hand, computer engineers and scientists do not know how to translate requirements for physical systems, such as **stability**, into computational requirements on performance, power consumption, etc."

"On the other hand, control and signal processing theory abstract computers largely as infallible numerical devices. This simplification **ignores** many important aspects of computing, such as increasingly larger **timing variance** due to caches and energy management and increasingly higher software error rates caused by complexity."

Again, a unified modeling environment supporting integrated simulation of all these aspects is needed. Finally, it is concluded that "Simplifying assumptions are also made about communications. Initial designs assume zero-loss, zero-delay communications, while neither occur in the wireless, low-power, shared, rapidly changing systems used in most CPS. The viability of future CPS must also address noise in measurements, inaccuracies in actuation, disturbances from the environment, and faults and failures in the computational process in a coherent, unified framework."

The motivation for this paper is to introduce the CPS field to the Modelica community and to describe how recent Modelica developments will facilitate integrated modeling and simulation of the cyber and physical parts of complex, hierarchical systems. The presented framework allows simulation of nonideal behavior, such as delays, noise, and quantization, as well as detailed timing simulations of realtime task scheduling and network communication. A robotics system will be used as demonstrator.

1.1 Related Work

The timing simulations presented in this paper are based on the TrueTime simulator [3, 4] developed at the Department of Automatic Control, Lund University, Sweden. The original version of TrueTime is based on MATLAB/Simulink [5] and the first Modelica-based version of TrueTime was presented in [6]. The version in [6], however, only dealt with the network simulation part of TrueTime, whereas the timing simulations presented in this paper also include simulation of real-time task scheduling in Modelica. In this paper, it is also shown how the timing simulations are integrated in a unified Modelica framework for systems configuration and simulation.

Most of the material presented in this paper was first prepared for a Modelica tutorial organized by the authors at the 3rd international CPSWeek multi-conference in Stockholm, April 2010.

1.2 Outline

The rest of this paper is outlined as follows. Section 2 continues the discussion on Modelica in the context of CPS and presents the recent development of Modelica for embedded systems with focus on features for simulation of non-ideal effects in controllers. Detailed simulation of controller timing variations in this new Modelica framework is introduced in Section 3. The cyber-physical simulation features are demonstrated on a robot example in Section 4. Finally, Section 5 gives the conclusions and directions for future work.

2 Modelica and CPS

The Cyber-Physical Systems Summit identified a series of applications that should give substantial societal impact while also presenting significant challenges. Three of these applications were developed into main grand challenges for CPS; *future distributed energy systems, future transportation systems*, and *next-generation healthcare systems*.

The first two challenges represent areas where the physical modeling capabilities of Modelica already are very strong, and medical and biological systems (third challenge) also contain physical elements well suited for equation-based modeling. Other applications of CPS include autonomous automotive systems, automatic pilot avionics, and distributed robotics.

There are, thus, many applications of CPS that require detailed physical modeling. However, since the physical modeling capabilities of Modelica are well established, this paper focuses on recent Modelica features that will allow modeling and cosimulation of the "cyber" aspects of CPS together with the physical models. The following aspects need to be addressed:

- The scale and complexity of CPS require compositional methods for integrated design and modeling.
- Most CPS applications are built up from controlled sub-systems, systems-of-systems, with local interaction between physical systems and controllers affecting global performance.
- Distributed sensing, actuation, and control needs to be modeled and simulated.
- The interfaces between the cyber and physical needs to be identified and properties of the interfaces should be easy to specify.
- New hybrid models may be needed to completely model all aspects of CPS.

2.1 Modelica for Embedded Systems

A first step towards accomplishing the aspects above in Modelica has been performed in the ITEA2 project EUROSYSLIB [7]. Language constructs were added to support embedded systems modeling and configuration using Modelica. A support library, *Modelica_EmbeddedSystems*, was developed as a user-friendly interface to the new language constructs. The new language constructs and a preliminary version of this library were presented in [8]. The effort to improve the cyber aspects of Modelica has continued in the ITEA2 MODELISAR [9] project, focusing on timing simulations and connections to the AUTOSAR [10] standard.

The new Modelica language features for embedded systems exhibit two important properties strongly related to the CPS aspects outlined above.

- 1. In complex models, constructed as a hierarchy of controlled sub-systems, a Modelica translator will be able to automatically deduce which parts that are truly physical, and which parts that are digital controllers ("cyber").
- 2. Special *communication blocks* are inserted in the interfaces between the cyber and physical parts. These contain replacable components, which makes it possible to easily define different properties (such as timing variations, noise, and inaccuracies) without modifying the structure of the original model.

2.2 Simulated Communication

The Modelica_EmbeddedSystems library contains special communication block implementations allowing simulation of non-ideal communication between controllers and plants. The effects that can be simulated are; sampling, computational delay (fraction of sample period), communication delay (unbounded), measurement noise, signal limitations, and quantization.

These effects are easily configured in the parameter dialog shown in Figure 1 below. The dialog is accessed from the communication blocks that make up the interfaces between controller and plant components of the controlled sub-systems.

| seneral Add modi | fiers | | | |
|------------------------|-----------------------------|----------|--|------------------------------------|
| omponent | | | | Icon |
| Name redeclare | fromInputToOutput | | | SimulatedRe |
| Comment | | | | write reade |
| odel | | | | |
| Path Modelica_E | EmbeddedSystems.Communicati | on.Simul | ated.SimulatedRealCommunication | |
| Comment Simulated I | DA or AD converter block | | | |
| ampling and noise — | | | | |
| sampled | | | = true, if output is sampled | |
| baseSampleRate | 0.005 | • 5 | Base sample rate of task | |
| sampleFactor | 1 | • | Sample factor relative to base sample period | bd |
| computationalDelay | 0 | • | Computational delay given as fraction of o | ne sample period (min=0,max=1) |
| communicationDelay | 0 | • 5 | Time delay due to communication | |
| noisy | false 🔽 | ٠ | = true, if output should be superimposed v | with noise |
| noise | 🕒 UniformNoise 🗸 🔝 | • | | |
| miting and quantizatio | in | | | |
| limited | V | | true, if output is | s limited |
| quantized | | | true 🔽 • 🛛 = true, if output o | uantization effects included |
| max | | | 50 • Upper limit of outp | ut (if limited = true) |
| min | | | -50 • Lower limit of outp | ut (if limited = true) |
| bits | | | 4 Number of bits of | quantization (if quantized = true) |
| | | | | |

Figure 1: Parameter dialog for the simulated communication block.

3 Timing Simulations

The simulated communication described above has been extended to enable more realistic simulation of timing effects. For example, latencies due to computation and/or communication are rarely constant but will vary as a result of the chosen task scheduling policy and communication protocol. These variations, called jitter, may also affect the actual sampling interval and could have a degrading effect on the performance of closed-loop control systems [3].

The developed timing simulations mimic the temporal behavior of real-time operating systems (RTOS) executing tasks in embedded processors and transmission of messages over communication networks. Aspects that can be simulated include latencies and jitter due to preemption and real-time task scheduling, queuing in the sending and receiving network nodes, latencies due to signal transmission, and latencies and jitter due to collisions and media access control (MAC) policies.

The simulations focus on control networks, i.e., networks sending relatively small packets regularly with tight real-time constraints [11]. The following real-time network protocols [12,13,14,15] are currently supported; Ethernet, CAN, TTCAN, FlexRay, and static time-triggered communication (TDMA).

3.1 Simulation Configuration

Task and subtask configuration records, provided in the Modelica_EmbeddedSystems library, are used to specify the real-time attributes (periods, offsets, priorities, and deadlines) of the simulated controller tasks. In addition, special kernel and network simulator models have been developed in Modelica to realize the timing simulators. These blocks use Modelica_EmbeddedSystems configuration records as input as shown for the kernel simulator in Figure 2.

| / | Config.kernetSimulator in ModelisarDemo.ECUTaskScheduting |
|--|---|
| | Comport Name Config.kernelSimulator Commert Loss Commert Lass Loss Loss Loss Loss Loss Loss Loss |
| suit tasi2 | Path Modelac_ImbeddedSystems: Configuration: TrueTime: RealTimeIternel Commet TrueTime: Interference Kennel Configuration Scheduling policy number: OfTasils 2. Jumber: OfTasils 2. KarRecords Subtask1, jubbask2, Jult |
| | CK Info Carcel |
| kornelSinukator Trus Tine Ke DEKEP O D D 2 tasks | |

Figure 2: Configuration of the kernel simulator.

The kernel simulator is configured by providing the desired scheduling policy, the number of simulated tasks, and an array of subtask records. The kernel simulator currently supports the following scheduling policies [16]:

- Fixed-priority scheduling
- Rate-monotonic scheduling
- Deadline-monotonic scheduling
- Earliest-deadline-first scheduling

The network configuration consists of general parameters (MAC policy, number of frames, data rate, minimum frame size, and loss probability) and policy-specific parameters. Policy-specific parameters needs to be specified for TDMA (Time Division Multiple Access), TTCAN, and FlexRay and consist of slot sizes and static communication schedules (represented as arrays of frame identifiers). In addition, FlexRay configuration also needs a communication schedule for its dynamic segment and a network idle time (in bits). The complete parameter dialog is shown in Figure 3.

| Seneral Add modif | iers |
|------------------------|--|
| omponent | Con |
| Name bus | RealTimeNet |
| Comment | |
| lodel | kbit/s |
| Path Modelica_E | mbeddedSystems.Configuration.TrueTime.RealTimeNetwork |
| Comment TrueTime ne | etwork simulator |
| eneral network configu | uration |
| MACPolicy | Jated.TrueTime.MacPolicyEnum.CAN 💙 • MAC policy for bus scheduling |
| numberOfFrames | 2 Number of frames |
| dataRate | 250 Data rate (kbit/s) |
| minimumFrameSize | 44 Minimum frame size (bits) |
| lossProbability | 0.0 • Loss probability |
| DMA configuration — | |
| slotSizeTDMA | 512 • Size of time slot (bits) |
| numberOfSlots | 3 Length of cyclic schedule (slots) |
| schedule | {1,1,2} F Cyclic schedule |
| FCAN configuration — | |
| slotSizeTTCAN | 512 • Size of time slot (bits) |
| numberOfCycles | 1 Number of base cycles |
| cycleLength | 2 • Length of base cycle (slots) |
| systemMatrix | [100, 1] System matrix, use ID = 100 for CAN arbitration |
| exRay configuration - | |
| staticSlotSize | 512 • Size of static time slot (bits) |
| numberOfStaticSlots | 3 Length of static schedule (slots) |
| staticSchedule | {1,1,2} III • Static schedule |
| miniSlotSize | 16 • Size of mini-slot (bits) |
| numberOfMiniSlots | Number of mini-slots in dynamic segment |
| dynamicSchedule | {1,1,2,2} Dynamic (mini-slot) schedule |
| networkIdleTime | 16 • Network idle time (bits) |
| | |

Figure 3: Configuration of the network simulator.

3.2 Communication

Having configured the kernel and/or network simulators, it is also required to configure individual signals in the model to use the timing simulators. This is done using the parameter dialog of the communication blocks. This configuration involves specifying the execution times of tasks, sizes and priorities of network frames, and references to the kernel and network simulators associated with the signal, see Figure 4.

| General Iruetime | Add modifiers | |
|------------------------|---------------------------------------|---|
| Kernel simulation | | |
| kernelSimulation | true 💌 🕨 | true if effects of task scheduling shall be simulated |
| executionTimeCalculate | executionTimeFactor*0.0005 > | Execution time to calculate outputs |
| executionTimeUpdate | executionTimeFactor*0.002 • | Execution time to update states |
| kernelHandle | config.kernelSimulator.kernelHandle > | Handle to kernel simulator object |
| Network simulation ——— | | |
| networkSimulation | true 💙 🔸 | true if effects of bus scheduling shall be simulated |
| frameDataSize | 601 • | Frame size (bits) |
| frameID | 1 • | Frame ID (1,,numberOfFrames) |
| framePriority | frameID | Priority of frame (for CAN and TTCAN) |
| networkHandle | bus.networkHandle + | Handle to network simulator object |
| | | |
| | | |
| | | |
| | | |

Figure 4: Configuration of individual signals.

3.3 Simulation Output

The kernel and network simulator models generate traces according to Figure 5. For each task in a simulated kernel, a task activation graph is produced. The graph has three levels, corresponding to idle, preempted by a higher-priority task, and running. Similarly for network frames, it is shown when a frame is being transmitted and when it is waiting because another transmission is using the bus. These graphs can be used to identify and verify timing properties.



Figure 5: Example of timing graph.

3.4 Modelica Implementation

The original TrueTime simulator is a Simulink Sfunction implemented in C++. The Modelica implementation uses the C external function interface to call the simulator at appropriate time instants.

The simulator is event-based and triggered based on internal and external events. Each time the simulated real-time kernel or network is executed, it computes the time for next internal, scheduled, event. The Modelica wrapper of the simulator then makes sure to generate a time event at the scheduled event. Examples of scheduled events are the completion of a network transmission or a real-time task finishing execution.

In between scheduled events, the kernel or network simulators could also be triggered by external events. External events include sending of a network message and triggering of task jobs.

3.5 Use Cases and Possible Extensions

The timing simulations can be used both to generate and to verify timing requirements of controller components. For example, with the new FMI [17] interface for model exchange and co-simulation it will be straight-forward to import a controller component, hook it up to the timing simulator and co-simulate it with its physical environment.

The current timing simulations assume that accurate execution time estimates are available for all tasks. A future extension will be to integrate the task simulations with tools for execution-time analysis. It will also be investigated how to integrate the timing simulations and task configurations with the timing model available in the latest AUTOSAR specification, version 4.0.

4 Simulation Examples

This section will demonstrate the simulation capabilities on a simple robotics example. It will first be shown how to simulate simple non-ideal behavior using the various options in the simulated communication parameter dialog of Figure 1. Then it will be demonstrated how to perform more detailed timing simulations including high-priority tasks and network transmissions.

4.1 Simulation Model

The simulation model, shown in Figure 6, contains three main components on the top level, a path planner, a controlled servo, and the mechanical model of a one-armed robot.



Figure 6: Simulation example.

The servo component contains a DC motor connected to a non-linear gear controlled in a three-level cascade with control loops for the motor angle, angular velocity, and motor current. Figure 7 shows the internals of the servo component. The controller component contains the angle and velocity P- and PIcontrollers, and it is shown how communication blocks have been inserted at the interfaces between the controllers and the physical plant.



Figure 7: Controlled servo extended with communication blocks.

4.2 Simulating Non-ideal Communication

The simulations below will consider a step change of ten degrees in the angle reference and the plots will show the angle, angular velocity and the controller output of the velocity loop (used as reference to the current controller). The ideal continuous case is included for reference in the simulation graphs.

We will first consider the effects of using sampled instead of continuous controllers. This is achieved by marking the checkbox *sampled* and selecting a sample rate (see Figure 1). In this case the sample rate is deliberately chosen too slow in relation to the closed-loop dynamics, and the results of that simulation are shown in Figure 8. The effects of too slow sampling are clearly seen in the velocity signal and in the control signal.



Figure 8: Simulation of sampled controller.

Next we will impair the stability properties of the system even more by also adding delay to the control signal computations. The results are shown in Figure 9, and in this case the performance of the system deteriorates to the point where the system starts to oscillate.



Figure 9: Sampled controller with delay.

Figure 10 shows the effects of adding measurement noise to the velocity measurement and it is seen how the noise is propagated and amplified by the velocity controller. Finally, Figure 11 demonstrates the possibility to simulate limitations on signals and signal quantization. In this case, the velocity controller output is limited between -50 and 50 with a very coarse resolution of only four bits.



Figure 10: Sampled controller with measurement noise.



Figure 11: Sampled controller with output quantization.

4.3 Timing Simulations

For the detailed timing simulations, we extend the robot model with a configuration component (including timing simulators and records to specify task attributes) and disturbing computations and communication according to Figure 12.



Figure 12: Extended robot model for timing simulation.

We will first examine the effects of real-time task scheduling. The simulated real-time kernel contains four tasks, representing the angle controller, velocity controller, current controller, and a disturbance task. The angle and velocity controller tasks have a period of 5 ms, whereas the current controller is executed with a period of 1 ms. The disturbance task runs on a 10 ms period. The simulated execution times are 100 μ s for the controller tasks and 1 ms for the disturbance task.

Figure 13 shows the result when the disturbance task is given the highest priority. In this case, the disturbance task interferes with the controller tasks every 10^{th} sample and the timing variations, as shown in the task schedule in Figure 14, causes the performance to deteriorate.



Figure 13: Simulation performance with highpriority disturbance task.



Figure 14: Task schedule with high-priority disturbance task.

By changing the scheduling policy to ratemonotonic scheduling (the shorter period, the higher priority), the controller task will automatically get higher priority than the slow disturbance task. In this case, the control performance is the same as for the 5 ms sampled simulation shown in Figure 8. The task schedule for this simulation is shown in Figure 15.



Figure 15: Task schedule when using rate-monotonic scheduling.

Finally, we will consider also network communication. In this simulation we assume that the control computer and the actuator reside on different nodes. The control signal computed by the inner controller (the current controller) then needs to be sent over the network. We also assume that the disturbance node sends high-priority traffic on the network.

We will compare two different network policies, CAN and FlexRay. CAN schedules frames according to priority, whereas FlexRay has a static segment where frames get exclusive network access during certain time slots. Figures 16 and 17 show the different transmission schedules obtained for CAN and FlexRay. For CAN, the disturbance frames interfere with the control signal transmissions. In the case of FlexRay, the controller frames are transmitted in the static segment and are then not affected by the disturbance frames.



Figure 16: Transmission schedule using CAN configuration.



Figure 17: Transmission schedule using FlexRay configuration.

5 Conclusions

This paper has introduced the field of cyber-physical systems (CPS) and has outlined some first steps in supporting CPS modeling and simulation using Modelica. It was shown how non-ideal controller behavior could be conveniently configured and simulated using the Modelica_EmbeddedSystems library. The simulated phenomena included sampling, delays, noise, limited signals, and quantization.

It was further demonstrated how detailed timing simulations based on the TrueTime simulator has been integrated in the Modelica_EmbeddedSystems framework. The presented timing simulations included both network simulations and simulations of task scheduling in real-time kernels. A robot example was used as a demonstrator of the presented simulation features.

5.1 Future Work

Future work will mainly focus on three areas; integration with tools for execution-time analysis, connections to the AUTOSAR 4.0 timing model, and extensions to support simulation of wireless network protocols. The first two areas were elaborated in Section 3.5.

Most CPS applications rely on wireless sensor networks in order to collect measurements that are used to close feedback loops. Thus, modeling of wireless network protocols and simulation of aspects related to this, such as signal attenuation, fading, packet losses, power consumption, etc is important in order to fully support CPS modeling and simulation. The Modelica version of TrueTime presented in [6] supports the wireless network protocols IEEE 802.11b/g (WLAN) and 802.15.4 (ZigBee).

5.2 Acknowledgements

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